

COMPUTING LOCALIZATIONS ITERATIVELY

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ABSTRACT. Let $R = \mathbb{C}[\mathbf{x}]$ be a polynomial ring with complex coefficients and $D_X = \mathbb{C}\langle \mathbf{x}, \boldsymbol{\partial} \rangle$ be the Weyl algebra. Describing the localization $R_f = R[f^{-1}]$ for nonzero $f \in R$ as a D_X -module amounts to computing the annihilator $A = \text{Ann}(f^a) \subset D_X$ of the cyclic generator f^a for a suitable negative integer a . We construct an iterative algorithm that uses truncated annihilators to build A for planar curves.

INTRODUCTION

Let $f \in R := \mathbb{C}[x_1, \dots, x_n]$ be a nonzero polynomial (n being a positive integer) and R_f the localization of the polynomial ring R with respect to f . Elements in R_f are rational functions $\frac{g}{f^k}$ with $g \in R$ and $k \in \mathbb{N}$. If f is not a constant, the R -module R_f is not finitely generated. One fundamental result by J. Bernstein states that R_f is finitely generated when considered as a left module over the complex Weyl algebra $A_n := A_n(\mathbb{C})$ of order n (see the needed definitions and precise statements in the next Section). This is a fundamental result in \mathcal{D} -module theory (the algebraic theory of systems of linear partial differential equations) which can be considered as part of Singularity Theory.

Bernstein's result states even more: the module R_f is cyclic over A_n and in fact it is generated by a rational function of the form f^a for some negative integer a . That means that R_f can be described as a quotient $\frac{A_n}{\text{Ann}(f^a)}$ where $\text{Ann}(f^a) := \text{Ann}_{A_n}(f^a)$ is the left ideal in A_n formed of the linear differential operators annihilating the rational function f^a . By [14] (see also [29] and [23]) one can take $a = -n + 1$.

One main problem in algorithmic \mathcal{D} -module theory is to compute a finite system of generators of the annihilating ideal $\text{Ann}(f^a)$. There are several algorithms solving this problem (see e.g. [19], [21], [20]),

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which use Gröbner bases and elimination theory in the Weyl algebra A_n . In the worst case, computing Gröbner bases has a doubly exponential complexity in both (commutative) polynomial rings and the Weyl algebra (see [1] and [13]); however, in practice, one should expect much longer running times for the latter on the input of the same size.

As the Weyl algebra A_n is a Noetherian ring, one can associate to the couple (f, a) the smallest integer $\kappa(f^a)$ such that $\text{Ann}(f^a)$ is generated by operators of order less than or equal to $\kappa(f^a)$. This numerical invariant plays a relevant role in the so called Logarithmic Comparison Problem with respect to the hypersurface defined in the complex affine space \mathbb{C}^n by the polynomial equation $f = 0$ (see e.g. [10], [27]).

In this paper we describe a new algorithm computing $\text{Ann}(f^{-1})$ for any reduced complex polynomial $f = f(x, y)$ in two variables. This algorithm uses Gröbner bases techniques in both polynomial rings and the Weyl algebra A_2 , but avoids elimination theory in the latter: our experiments show that the bottleneck of the algorithm is a syzygy module computation over the former.

We first compute iteratively, for $d = 1, 2, \dots$, truncated annihilating ideals $\text{Ann}^{(d)}(f^{-1})$ generated by linear differential operators annihilating f^{-1} and of order less than or equal to d . The ideal $\text{Ann}^{(d)}(f^{-1})$ can be computed by using polynomial Gröbner basis. We then compute equations for the characteristic cycle of the ideal $\text{Ann}^{(d)}(f^{-1})$ and compare it with the one of $\text{Ann}(f^{-1})$ which only depends on the multiplicity of the plane curve $f = 0$ at its singular points; for the beforementioned comparison we only need to localize at each of these singular points.

The paper is organized as follows. In the first Section we survey basic results on the Weyl algebra, holonomic D -modules, b -functions (or Bernstein-Sato polynomials) and the annihilating ideal of some rational functions on \mathbb{C}^n . In Section 2 we describe the new iterative algorithm, we give a stopping criterion and prove its correctness. We also treat some examples, in particular the so called family of Reiffen's curves $f_{p,q}$. Our implementation of the algorithm in Macaulay2 works on larger examples of Reiffen's curve family than the known general algorithm and our experiment demonstrates an interesting (very simple!) dependence of the order of generation of $\text{Ann}^{(d)}(f_{p,q}^{-1})$ on p . In the last Section we conclude the discussion and propose some open questions on the subject.

1. PRELIMINARIES

1.1. Weyl algebra. Let $n \geq 0$ be an integer and $X = \mathbb{C}^n$ be the n -dimensional complex affine space. Define the n -th *Weyl algebra* as the

associative algebra

$$D_X = \mathbb{C}\langle \mathbf{x}, \boldsymbol{\partial} \rangle = \mathbb{C}\langle x_1, \dots, x_n, \partial_1, \dots, \partial_n \rangle$$

where $[\partial_i, x_i] = \partial_i x_i - x_i \partial_i = 1$ and all other pairs of generators commute. The Weyl algebra D_X is isomorphic to the algebra of *linear differential operators* with coefficients in the polynomial ring $R = \mathbb{C}[\mathbf{x}] = \mathbb{C}[x_1, \dots, x_n]$. Every element in D_X has a unique *normal form*

$$Q = \sum_{\alpha, \beta \in \mathbb{N}^n} q_{\alpha\beta} \mathbf{x}^\alpha \boldsymbol{\partial}^\beta,$$

where finitely many of $q_{\alpha\beta} \in \mathbb{C}$ are nonzero and where $\mathbf{x}^\alpha \boldsymbol{\partial}^\beta$ stands for $x_1^{\alpha_1} \dots x_n^{\alpha_n} \partial_1^{\beta_1} \dots \partial_n^{\beta_n}$. We denote the n -th Weyl algebra A_n if we need to emphasize the dimension of X ; throughout this paper $D_X = A_2$ is used the most.

The ring D_X is simple: there are only trivial two-sided ideals. All D_X -ideals and D_X -modules considered in this article are *left* ideals and modules, respectively.

Examples of D_X -modules include functional spaces, e.g., polynomial functions $\mathbb{C}[\mathbf{x}]$ and smooth functions $C^\infty(X)$, and power series rings $\mathbb{C}[[\mathbf{x}]]$.

Another example is the localization of the polynomial ring $\mathbb{C}[\mathbf{x}, f^{-1}]$ where f is a nonzero polynomial with the natural action defined as follows:

$$\begin{aligned} x_i \cdot g f^{-j} &= x_i g f^{-j}, \\ \partial_i \cdot g f^{-j} &= \left(\frac{\partial g}{\partial x_i} f - j g \frac{\partial f}{\partial x_i} \right) f^{-j-1}, \end{aligned}$$

for $1 \leq i \leq n$, $g \in \mathbb{C}[\mathbf{x}]$, and $j \in \mathbb{N}$.

1.2. Gröbner bases. It is possible to compute in the Weyl algebra D_X , since it is *Gröbner friendly*; see e.g. [8, 9], [24] and [15].

Gröbner bases can be computed with respect to any w -compatible monomial order, where $w = (w_{\mathbf{x}}, w_{\boldsymbol{\partial}}) \in \mathbb{R}^{2n}$ satisfies $w_{\mathbf{x}} + w_{\boldsymbol{\partial}} \geq \mathbf{0}$ componentwise. This condition, in particular, guarantees that the filtration $\{F_i^w\}_{i \in \mathbb{Z}}$ of D_X , where

$$F_i^w = \text{Span} \{ \mathbf{x}^\alpha \boldsymbol{\partial}^\beta : \alpha, \beta \in \mathbb{N}^n, w \cdot (\alpha, \beta) \leq i \},$$

is preserved under taking the commutator. Therefore, the associated graded algebra $\text{gr}_w(D_X)$ is well defined. Note that

- if $w_{\mathbf{x}} + w_{\boldsymbol{\partial}} = \mathbf{0}$ then $\text{gr}_w(D_X) = D_X$;

- if $w_{\mathbf{x}} + w_{\partial} > \mathbf{0}$ componentwise then

$$\mathrm{gr}_w(D_X) \cong \mathbb{C}[\mathbf{x}, \boldsymbol{\xi}] = \mathbb{C}[x_1, \dots, x_n, \xi_1, \dots, \xi_n],$$

a polynomial ring in $2n$ variables.

The w -order of a nonzero operator $Q = \sum_{\alpha, \beta \in \mathbb{N}^n} q_{\alpha\beta} \mathbf{x}^\alpha \partial^\beta \in D_X$, denoted $\mathrm{ord}_w(Q)$, is defined as the maximum of $w \cdot (\alpha, \beta)$ for $q_{\alpha\beta} \neq 0$. The initial part of Q with respect to w is

$$\mathrm{in}_w(Q) = \sum_{w \cdot (\alpha, \beta) = \mathrm{ord}_w(Q)} q_{\alpha\beta} \mathbf{x}^\alpha \boldsymbol{\xi}^\beta \in \mathrm{gr}_w(D_X).$$

We simply denote $\mathrm{in}_w(0) = 0$.

Definition 1. Let I be an ideal in D_X . The characteristic ideal

$$\mathrm{gr}_w(I) = \mathrm{in}_w(I) \subset \mathrm{gr}_w(D_X) = \mathbb{C}[\mathbf{x}, \boldsymbol{\xi}],$$

with $w = (0, e)$, i.e.: $w(x_i) = 0$ and $w(\partial_i) = 1$ for $i = 1, \dots, n$ defines the characteristic variety

$$\mathrm{Char}(D_X/I) = \mathbb{V}(\mathrm{gr}_w(I)) \subset \mathbb{C}^{2n}.$$

Here $\mathrm{in}_w(I)$ stands for the ideal of $\mathrm{gr}_w(D_X)$ generated by $\{\mathrm{in}_w(Q), | Q \in I\}$, the set of the leading forms of the elements of I with respect to the weight w . The set $\mathbb{V}(J)$, for an ideal J in $\mathbb{C}[\mathbf{x}, \boldsymbol{\xi}]$, is the variety in \mathbb{C}^{2n} defined by the ideal J .

If M is a nonzero finitely generated D_X -module, one can define, following [2], its characteristic variety $\mathrm{Char}(M) \subset \mathbb{C}^{2n}$ by using any *good filtration* on M . The set $\mathrm{Char}(M)$ is an affine algebraic subset of \mathbb{C}^{2n} . In particular one can consider its *Krull dimension* $\dim(\mathrm{Char}(M))$.

To each irreducible component C of $\mathrm{Char}(M)$ one can associate its *multiplicity* m_C which is a positive integer number. For the definition of the Krull dimension of an affine algebraic set and the multiplicity of its irreducible components one can consult [4].

Theorem 1 (Bernstein's inequality; [2]). *Let M be a nonzero, finitely generated D_X -module. Then $n \leq \dim(\mathrm{Char}(M)) \leq 2n$. In particular, if $I \subsetneq D_X$ is a left ideal, then $n \leq \dim(\mathrm{Char}(D_X/I)) \leq 2n$.*

There is a number of software implementations of D -module algorithms: `kan/sm1` [26], `risa/asir` [18], `dmod.lib` library [16] in Singular [11], and the `Dmodules` package [17] in Macaulay2 [12]. The last package hosts our implementations of the algorithms in this article.

1.3. Holonomic D -modules. A finitely generated D_X -module M is called *holonomic* if either $M = (0)$ or the dimension of its characteristic variety equals n (i.e. $\dim(M) = n$). An ideal I in D_X is called holonomic if either $I = D_X$ or D_X/I is holonomic.

The characteristic cycle of a holonomic D_X -module M is by definition the sum

$$\text{CChar}(M) := \sum_{C \text{ irreducible component of } \text{Char}(M)} m_C C$$

viewed as a cycle in the cotangent bundle $T^*\mathbb{C}^n$. Here m_C is the multiplicity of C in $\text{Char}(M)$.

Holonomic modules are particularly nice from the point of view of computations due to the following fact.

Theorem 2 (Stafford [25]). *Every holonomic D_X -module is cyclic.*

Every holonomic module M can be thought of as $M = D_X \cdot Q$, where Q is a cyclic generator, hence, $M \cong D_X / \text{Ann}_{D_X}(Q)$.

One example of a holonomic module is $R = \mathbb{C}[\mathbf{x}]$; it is generated by $1 \in R$. Another is its localization $R_f = \mathbb{C}[\mathbf{x}, f^{-1}]$, as stated by a theorem by J. Bernstein.

Theorem 3 ([2, 3]). *Let $f \in \mathbb{C}[\mathbf{x}]$ be a nonzero polynomial. Then the D_X -module R_f is holonomic. Moreover, for some negative integer a the element $f^a \in R_f$ generates R_f as a D_X -module.*

Computing the annihilator $\text{Ann}_{D_X}(f^a)$ is the main topic of this article.

By [2, 3] the largest possible exponent $a = a(f)$ above is equal to the smallest integer root of the *Bernstein-Sato polynomial* $b_f(s) \in \mathbb{Q}[s]$, defined as the monic nonzero polynomial $b(s)$ of the smallest possible degree such that

$$b(s)f^s = Q(s, \mathbf{x}, \boldsymbol{\partial}) \cdot f^{s+1}, \text{ where } Q \in D_X[s] := \mathbb{C}[s] \otimes_{\mathbb{C}} D_X.$$

The idea of many localization algorithms (first formulated by Oaku [19] in 1996) is simple to state:

- (1) Compute $\text{Ann}(f^s)$ as an ideal of $D_X[s]$.
- (2) Compute $b_f(s)$ to determine $a = a(f)$.
- (3) Specialization: $\text{Ann}(f^s)|_{s=a}$.

This approach, for example, is implemented in the function `Dlocalize` of the `Dmodules` package. The bottlenecks of the algorithm are items 1 and 2 that require an expensive elimination via Gröbner bases in the Weyl algebra D_X .

Step 2 can be avoided by using the following estimate

Theorem 4 ([14], see also [29], [23]). R_f is generated by f^{-n+1} .

As a corollary, for planar curves $f \in \mathbb{C}[x, y]$, localization $R_f = D_X \cdot f^{-1} = D_X / \text{Ann}(f^{-1})$.

Remark 1. Note that one can easily compute the annihilator of a rational function in the algebra of linear differential operators with rational function coefficients $\mathbb{C}(\mathbf{x})[\partial] = \mathbb{C}(\mathbf{x}) \otimes_{\mathbb{C}[\mathbf{x}]} D_X$: the annihilator is a maximal ideal of $\mathbb{C}(\mathbf{x})[\partial]$, e.g.,

$$\text{Ann}_{\mathbb{C}(\mathbf{x})[\partial]}(f^a) = \mathbb{C}(\mathbf{x}) \otimes \text{Ann}_{D_X}(f^a) = \langle f\partial_i - a\frac{\partial f}{\partial x_i} \mid i = 1, \dots, n \rangle.$$

2. ITERATIVE ALGORITHM

Our main example will be the family of the so-called Reiffen's curves; see [22].

$$f = f_{p,q} = x^p + y^q + xy^{q-1} = 0,$$

where $p \geq 4, q \geq p + 1$.

Example 1. Computing the annihilator (this can be done via `AnnFs` followed by specialization $s = -1$),

$$\begin{aligned} \text{Ann}(f_{4,5}^{-1}) = & \langle 4x^2\partial_x + 5xy\partial_x + 3xy\partial_y + 4y^2\partial_y + 16x + 20y, \\ & 16xy^2\partial_x + 4y^3\partial_x + 12y^3\partial_y - 125xy\partial_x - 4x^2\partial_y + \\ & 5xy\partial_y - 100y^2\partial_y + 64y^2 - 500y, \\ & 16y^3\partial_x^2 - 16y^3\partial_x\partial_y + 125xy\partial_x^2 - 35xy\partial_x\partial_y + 100y^2\partial_x\partial_y + \\ & 12x^2\partial_y^2 - 2xy\partial_y^2 - 24y^2\partial_y^2 + 112xy\partial_x - 36y^2\partial_x + \\ & 84y^2\partial_y - 930x\partial_x + 625y\partial_x + 26x\partial_y - \\ & 893y\partial_y + 448y - 3720 \rangle, \end{aligned}$$

we see that it is generated in order 2, i.e., generated by operators of order at most 2.

One interesting question is: What is the order of generation of this annihilator for a given p and q ?

2.1. Iterative approach.

Definition 2. We call

$$\text{Ann}^{(d)}(f^a) = \langle Q \in \text{Ann}(f^a) \mid \text{ord } Q \leq d \rangle$$

the d -th truncated annihilator of f^a .

To compute $\text{Ann}^{(d)}(f^a)$ one may find the R -syzygy module S_d for the vector of partial derivatives $(\partial^\alpha \cdot f^a)_\alpha$ where $|\alpha| \leq d$.

Example 2. To find an element in $\text{Ann}^{(1)}(f^{-1})$ for $f = x^2 - y^3$ consider all partial derivatives of f^{-1} of order at most 1:

$$\begin{aligned} \partial_x \cdot f^{-1} &= -2x f^{-2} \\ \partial_y \cdot f^{-1} &= 3y^2 f^{-2} \\ 1 \cdot f^{-1} &= (x^2 - y^3) f^{-2} \end{aligned}$$

For instance, $3x(-2x) + 2y(3y^3) + 6(x^2 - y^3) = 0$, hence,

$$3x\partial_x + 2y\partial_y + 6 \in \text{Ann}^{(1)}(f^{-1}).$$

In `Dmodules` the function for computing the truncated annihilator is called `kOrderAnnFa`:

```
i1 : loadPackage "Dmodules";
i2 : R = QQ[x,y];
i3 : f = x^2-y^3;
i4 : A1 = kOrderAnnFa(1,f,-1)

o4 = ideal (3x*dx + 2y*dy + 6, 3y^2 dx + 2x*dy, y^3 dy - x^2 dy + 3y^2)
```

Note that `kOrderAnnFa(d,f,-1)` would return the same ideal for all $d \in \mathbb{N}$ as the annihilator for this particular example is generated in order 1.

Remark 2. Suppose S_{d-1} , or in other words $\text{Ann}^{(d-1)}(f^a)$, is known. Then the computation of S_d , or the d -th truncated annihilator, can be optimized by computing the syzygies modulo S_{d-1} .

In practice, however, it appears that the time of the syzygy computation at order d dominates that for order $d-1$ for all d . In particular, in Experiment 1 for the Reiffen curve of any degree the last step takes more time than all previous steps combined.

2.2. Stopping criterion. The sequence of truncated annihilators stabilizes at some point:

$$\text{Ann}^{(1)}(f^a) \subseteq \text{Ann}^{(2)}(f^a) \subseteq \cdots \subseteq \text{Ann}^{(d)}(f^a) = \text{Ann}^{(d+1)}(f^a) = \cdots$$

Definition 3. The smallest d such that $\text{Ann}^{(d)}(f^a) = \text{Ann}(f^a)$ is denoted by $\kappa(f^a)$ and is called the annihilator order of f^a .

Note that $\text{Ann}^{(d)}(f^a) = \text{Ann}^{(d-1)}(f^a)$ does not imply $\text{Ann}^{(d+1)}(f^a) = \text{Ann}^{(d)}(f^a)$. For example, for $f = x$ and $a = 3$ we have $\text{Ann}^{(1)} = \text{Ann}^{(2)} = \text{Ann}^{(3)} = \langle x\partial_x - 3 \rangle$ and $\text{Ann}^{(4)} = \langle x\partial_x - 3, \partial_x^4 \rangle$.

2.3. Annihilator order of a planar curve. In view of Theorem 4, for a planar curve defined by $f \in R := \mathbb{C}[x, y]$ the localization R_f is generated by f^{-1} . We have constructed an algorithm that computes truncated annihilators stopping exactly at the annihilator order of f^{-1}

thus producing the whole annihilator $\text{Ann}(f^{-1})$ for a plane curve with at most one singular point, that we can assume to be the origin.

Algorithm 1. $(d, A) = \text{KAPPAANDANNIHILATOR}(f)$

Input: $f \in R = \mathbb{C}[x, y]$, a plane curve with at most one singular point, that we can assume to be the origin.

Output: $d = \kappa(f^{-1})$, $A = \text{Ann}(f^{-1})$.

1: $d \leftarrow 0$.

2: **repeat**

3: $d \leftarrow d + 1$.

4: $A \leftarrow \text{Ann}^{(d)}(f^{-1})$.

5: $\mathfrak{a} \leftarrow$ primary component of $\text{gr}_w(A)$ corresponding to the origin.

6: **until** $\deg \mathfrak{a} = (\text{multiplicity of } f \text{ at the origin}) - 1$

Proof of correctness. We see $X = \mathbb{C}^2$ as a complex manifold and consider C , the analytic plane curve defined by the polynomial equation $f = 0$. We assume f to be a reduced polynomial. Denote

- by \mathcal{O}_X the sheaf of holomorphic functions on X ,
- by \mathcal{D}_X the sheaf of linear partial differential operators on X with holomorphic coefficients, and
- by $\mathcal{O}_X[*C]$ the sheaf of meromorphic functions on X with poles on C .

By [14] the sheaf $\mathcal{O}_X[*C]$ is a coherent holonomic \mathcal{D}_X -module. If $p \in X \setminus C$ then $\mathcal{O}_X[*C]_p$ is just $\mathcal{O}_{X,p}$ and if the point $p \in C$ is a smooth point then the $\mathcal{D}_{X,p}$ -module $\mathcal{O}_X[*C]_p$ is isomorphic to the quotient

$$\frac{\mathcal{D}_{X,(0,0)}}{\mathcal{D}_{X,(0,0)}(x\partial_x + 1, \partial_y)}$$

(for some choice of local coordinates (x, y)). So, in the neighborhood of a smooth point in C , the characteristic variety of $\mathcal{O}_X[*C]$ is just the union of the conormal space to C , say T_C^*X and the zero section T_X^*X . Moreover, the multiplicities of each of these components in the corresponding characteristic cycle is 1.

It is well known that (see e.g. [5, Section 6]) in a sufficiently small neighborhood U of a point $p \in C$, the characteristic cycle of $\mathcal{O}_U[*C]$ is

$$\text{CChar}(\mathcal{O}_U[*C]) = T_U^*U + \overline{T_{U \cap C \setminus \{p\}}^*U} + (m - 1)T_p^*U$$

where T_U^*U is the zero section of the cotangent bundle T^*U , $\overline{T_{U \cap C \setminus \{p\}}^*U}$ stands for the closure in T^*U of the conormal bundle to the smooth part of $U \cap C$, T_p^*U is the conormal bundle to the point p and m is the multiplicity of the plane curve C at the point p .

Let us denote $\mathcal{I} := \mathcal{A}nn_{\mathcal{D}_U}(f^{-1})$ and for any integer number $d \geq 1$, $\mathcal{I}^{(d)} := \mathcal{A}nn_{\mathcal{D}_U}^{(d)}(f^{-1})$ considered as a sheaf of ideals in \mathcal{D}_U .

As f is a global equation for the plane curve $C \subset X$ we have $\mathcal{I}_q^{(d)} = \mathcal{D}_{U,q} \mathcal{A}nn_{\mathcal{D}_X}(f^{-1})$ for any $q \in U$.

We have the following exact sequence of holonomic \mathcal{D}_U -modules:

$$0 \rightarrow \frac{\mathcal{I}}{\mathcal{I}^{(d)}} \rightarrow \frac{\mathcal{D}_U}{\mathcal{I}^{(d)}} \rightarrow \frac{\mathcal{D}_U}{\mathcal{I}} \rightarrow 0$$

and then the following equality of characteristic cycles:

$$\text{CChar} \left(\frac{\mathcal{D}_U}{\mathcal{I}^{(d)}} \right) = \text{CChar} \left(\frac{\mathcal{D}_U}{\mathcal{I}} \right) + \text{CChar} \left(\frac{\mathcal{I}}{\mathcal{I}^{(d)}} \right).$$

Denote by $m^{(d)}$ the multiplicity of the irreducible component T_p^*U in $\text{CChar} \left(\frac{\mathcal{D}_U}{\mathcal{I}^{(d)}} \right)$. As the ideals $\mathcal{I}^{(d)}$ and \mathcal{I} coincide on $U \cap C \setminus \{p\}$ the holonomic \mathcal{D}_U -module $\frac{\mathcal{I}}{\mathcal{I}^{(d)}}$ is concentrated on p . Then $\text{CChar} \left(\frac{\mathcal{D}_U}{\mathcal{I}^{(d)}} \right) = \text{CChar} \left(\frac{\mathcal{D}_U}{\mathcal{I}} \right)$ if and only $m^{(d)} = m - 1$ (where m is the multiplicity of the plane curve C at $p \in C$). We also have $\mathcal{I}^{(d)} = \mathcal{I}$ if and only if $m^{(d)} = m - 1$.

Once we have a finite set of polynomial equations for the characteristic ideal $\text{gr}_w(\mathcal{A}nn^{(d)}(f^{-1}))$, these equations define the characteristic cycle $\text{CChar} \left(\frac{\mathcal{D}_U}{\mathcal{I}^{(d)}} \right)$ in the neighborhood U of any point $p \in X$. The multiplicity $m^{(d)}$ equals the degree of the primary ideal in the primary decomposition of $\text{gr}_w(\mathcal{A}nn^{(d)}(f^{-1}))$ corresponding to T_p^*U . \square

Remark 3. *The multiplicity $m^{(d)}$ can be also obtained as the multiplicity of the (unique) point $\tilde{p} \in T_p^*U$ in the intersection $\text{gr}_w(\mathcal{A}nn^{(d)}(f^{-1})) \cap L$ where L is the defining ideal of a generic 2-plane in the 4-dimensional ambient space.*

Remark 3 allows us to avoid the expensive primary decomposition step in the Algorithm 1. Assume $p = 0$ is the only singularity; hence, let $U = X$ and let $S = \mathbb{C}[x, y, \xi, \eta]$ be the ring of the characteristic ideal $\text{gr}_w(\mathcal{A}nn^{(d)}(f^{-1}))$. In our implementation we set $L = \langle \xi - a, \eta - b \rangle$, $(a, b) \in \mathbb{C}^2 \setminus \{(0, 0)\}$; in practice, a and b are taken to be small integers. The *genericity* of the plane L means that the point $\tilde{p} = (0, 0, a, b)$ does not belong to $\overline{T_{U \cap C \setminus \{p\}}^*U}$. The following algorithm checks the genericity of the choice.

Algorithm 2. $g = \text{CHECKGENERICITY}(f, a, b)$

Input: $f \in R = \mathbb{C}[x, y]$, a plane curve with the unique singular point at the origin.

Output: g , a boolean value, whether the genericity condition described in Remark 3 is satisfied.

- 1: Compute the syzygies $M \leftarrow \{(u, v) \in R^2 \mid uf_x + vf_y = 0\}$.
- 2: $I = \langle u\xi + v\eta \mid (u, v) \in S \rangle + \langle f \rangle$.
- 3: Compute the saturation $J \leftarrow I : \langle x, y \rangle^\infty$.
- 4: $g \leftarrow (0, 0, a, b) \in \mathbb{V}(J)$

Remark 4. The choice of the plane $L = \langle \xi - a, \eta - b \rangle$, $(a, b) \in \mathbb{C}^2 \setminus \{(0, 0)\}$, is artificial but sufficient, since $\dim(T_p^*U \cap T_{U \cap C \setminus \{p\}}^*U) \leq 1$. In fact, in Experiment 1 the choice $(a, b) = (0, 1)$ always works.

Remark 5. In view of Remark 3, $m^{(d)}$ can be computed numerically: one can pick a random L and compute an approximation to the point of intersection \tilde{p} . Using this approximation, one may recover the multiplicity $m^{(d)}$ heuristically. To the best of our knowledge there is no known approach to certification of the correctness of such a computation.

Moreover, in our experiments the bottleneck is not the saturation step in Algorithm 2, but the computation of syzygies (described after Definition 2).

Remark 6. The algorithm can be generalized for arbitrary complex planar curves C defined by a reduced polynomial f . As the number of singular points of C is finite it is enough to apply our iterative algorithm at each singular point of C and compute $\kappa(f_p^{-1})$ for each $p \in \text{Sing}(C)$. We finally have $\kappa(f^{-1})$ as the maximum of the $\kappa(f_p^{-1})$ for $p \in \text{Sing}(C)$. See also Section 3.1 for a generalization of this algorithm to a more general situation.

We have implemented Algorithm 1 modified according to the Remark 3 and including the genericity check provided by Algorithm 2 in the `Dmodules` package for Macaulay2. The following example demonstrates that the new function `kappaAnnF1PlanarCurve` produces the same result as the old method involving `AnnFs`.

```
i2 : f = reiffen(7,8)
o2 = x17x28 + x28 + x17
o2 : QQ[x1, x2]
i3 : As = AnnFs f;
o3 : Ideal of QQ[x1, x2, dx1, dx2, s]
i4 : A = sub(As, {last gens ring As => -1});
i5 : (kappa,A') = kappaAnnF1PlanarCurve f;
i6 : kappa
o6 : 4
```

```
i7 : A == sub(A', ring A)
o7 = true
```

Experiment 1. *We have computed the annihilator order for the Reiffen curve: $\kappa(f_{p,q}^{-1})$ seems to depend only on p . Starting with $p = 4$ (the last computed value is for $p = 21$), the sequence $\kappa(f_{p,*}^{-1})$ is*

$$2, 2, 3, 4, 4, 5, 6, 6, 7, 8, 8, 9, 10, 10, 11, 12, 12, 13, \dots$$

We can compute $\kappa(f_{p,q}^{-1})$ for $q \leq 22$ in less than one day using our implementation of Algorithm 1. In contrast, an implementation of the general algorithm (the computation of $\text{Ann}(f^s)$ followed by the specialization $s = -1$) took about a week for $p = 17$ and has not produced a result in one month for $p = 18$.

Another interesting observation is that the sequence of multiplicities $(m^{(d)})$ for $d = 1, \dots, \kappa(f_{p,q}^{-1})$ does not depend on q . Below we report the computed values $m^{(d)}$:

$d =$	1	2	3	4	5	6	7	8	9	10	11	12	13
$p = 4$	4	3											
$p = 5$	6	4											
$p = 6$	8	6	5										
$p = 7$	10	8	7	6									
$p = 8$	12	10	9	7									
$p = 9$	14	12	11	9	8								
$p = 10$	16	14	13	11	10	9							
$p = 11$	18	16	15	13	12	10							
$p = 12$	20	18	17	15	14	12	11						
$p = 13$	22	20	19	17	16	14	13	12					
$p = 14$	24	22	21	19	18	16	15	13					
$p = 15$	26	24	23	21	20	18	17	15	14				
$p = 16$	28	26	25	23	22	20	19	17	16	15			
$p = 17$	30	28	27	25	24	22	21	19	18	16			
$p = 18$	32	30	29	27	26	24	23	21	20	18	17		
$p = 19$	34	32	31	29	28	26	25	23	22	20	19	18	
$p = 20$	36	34	33	31	30	28	27	25	24	22	21	19	
$p = 21$	38	36	35	33	32	30	29	27	26	24	23	21	20

3. DISCUSSION AND OPEN PROBLEMS

We have proposed a new algorithm computing a presentation, as a left A_2 -module, of the localization $\mathbb{C}[x, y]_f$ for nonzero polynomials

$f = f(x, y)$ in two variables. The proof of the correctness of the algorithm is specific to dimension 2 and it seems difficult to generalize it to higher dimensions. Nevertheless, in the next subsection we describe an analogous algorithm for polynomials $f = f(x_1, \dots, x_n)$, for any $n > 0$, under the hypothesis that the hypersurface defined by f in \mathbb{C}^n has only isolated singularities.

3.1. Isolated hypersurface singularities. The iterative algorithm described before can be also applied in a more general situation. Let $n > 0$ be an integer. Denote $X = \mathbb{C}^n$, by \mathcal{O}_X the sheaf of rings of holomorphic functions on X and by \mathcal{D}_X the sheaf of rings of linear differential operators with holomorphic coefficients. Assume that $f \in R = \mathbb{C}[x_1, \dots, x_n]$ is a non zero reduced polynomial defining a hypersurface $C \in X$ with an isolated singularity $p \in C$. Then the characteristic cycle of $\mathcal{O}_X[*C]$ has the following description (see e.g. [5, Section 6])

$$\text{CChar}(\mathcal{O}_X[*C]) = T_X^*X + \overline{T_{C \setminus \{p\}}^*X} + \nu T_p^*X$$

where ν is the Milnor number of a generic hyperplane section of C through the point p . This formula coincides with the one we have considered before for $n = 2$. The number ν can be computed in an effective way. We can repeat the same ideas as in the proof of the correctness of Algorithm 1. By comparison of the multiplicity $m^{(d)}$ with ν , where $m^{(d)}$ is the multiplicity of the irreducible component T_p^*X in

$$\text{CChar}\left(\frac{\mathcal{D}_X}{\mathcal{D}_X \text{Ann}_{A_n}(f^{-n+1})}\right),$$

we can compute the annihilator order of f^{-n+1} .

3.2. Weyl closure.

Definition 4. *The Weyl closure of $I \subset D_X$ is defined as*

$$\text{WeylCl}(I) = K(\mathbf{x}) \otimes_{\mathbb{C}[\mathbf{x}]} I \cap D_X.$$

Note that the problem of computing $\text{Ann}(f^a)$ is a particular instance of the problem of computing the Weyl closure: according to 1 the annihilator of f^a in $K(\mathbf{x}) \otimes_{\mathbb{C}[\mathbf{x}]} D_X$ is equal to $\langle f\partial_i + \frac{\partial f}{\partial x_i} \mid i = 1, \dots, n \rangle$. Therefore,

$$\text{Ann}(f^{-1}) = \text{WeylCl}(\langle f\partial_i + \frac{\partial f}{\partial x_i} \mid i = 1, \dots, n \rangle).$$

There exists an algorithm due to Tsai [28] that computes the Weyl closure. However, the computation is highly nontrivial and boils down

to computing localizations of D_X -modules anyway. On the other hand, one may define the truncated Weyl closure of order d analogously to the d -th truncated annihilator. Then, the sequence of truncated closures stabilizes at some degree; finding this degree or discovering a termination criterion for an iterative algorithm similar to Algorithm 1 would result in a much simpler procedure.

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